



# Four Attacks and a Proof for Telegram

RWC 2022

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**Martin R. Albrecht, Lenka Mareková, Kenneth G. Paterson, Igors Stepanovs**

Royal Holloway, University of London; ETH Zürich







Based on the paper to appear at IEEE S&P 2022.

More information at: <https://mtpsym.github.io/>

# Background

## Monthly active users in Jan 2022:

*According to Statista 2022.*

 <b>WhatsApp</b>	$2000 \cdot 10^6$
 <b>WeChat</b>	$1263 \cdot 10^6$
 <b>FB Messenger</b>	$988 \cdot 10^6$
 <b>QQ</b>	$574 \cdot 10^6$
 <b>Snapchat</b>	$557 \cdot 10^6$
 <b>Telegram</b>	$550 \cdot 10^6$

## Collective Information Security in Large-Scale Urban Protests: the Case of Hong Kong

Martin R. Albrecht, Jorge Blasco, Rikke Bjerg Jensen, and Lenka Mareková, *Royal Holloway, University of London*

**Telegram** was the predominant messaging application used in the Hong Kong protests in 2019-2020.

**Telegram** was perceived to provide more security than its competitors.

Important advantages of **Telegram**:

- Support of public and private **group chats** for up to 200'000 people (**Signal** up to 1'000; **WhatsApp** up to 256).
- **Pseudonymity**: can use a pseudonym, not revealing phone number to others (not supported in **Signal** and **WhatsApp**).
- **Other features**: anonymous polls; disappearing messages; timed or scheduled messages; ability to delete messages sent by others.

Common use cases: large **public groups** up to 50'000 members, and small **private groups**.

# Cloud Chats and Secret Chats

	Cloud Chats	Secret Chats
Group communication	✓	✗
1-on-1 communication	✓	✓
Type of encryption	client-server	end-to-end
Enabled by default?	✓	✗

“Q: Why are you not using X? (insert solution)

While other ways of achieving the same cryptographic goals, undoubtedly, exist, we feel that the present solution is both robust and also succeeds at our secondary task of beating unencrypted messengers in terms of delivery time and stability.”

**Telegram FAQ** (<https://core.telegram.org/techfaq>)

**Why not use TLS instead of MTProto?**



The **MTProto** protocol – **Telegram’s** equivalent of the **TLS record protocol**.

**Cloud Chats** encrypt and authenticated messages using **MTProto**.

**Secret Chats** add another layer of **MTProto** encryption, i.e. messages are **double-encrypted**.

The **MTProto** protocol is not well-studied:

**2013: Telegram launched** with MTProto 1.0.

**2016: Jakobsen and Orlandi** showed that MTProto 1.0 is not CCA-secure.

**2017: Telegram released MTProto 2.0** that addressed the security concerns.

**2017: Sušánka and Kokeš** reported an attack based on improper validation in the Android client.

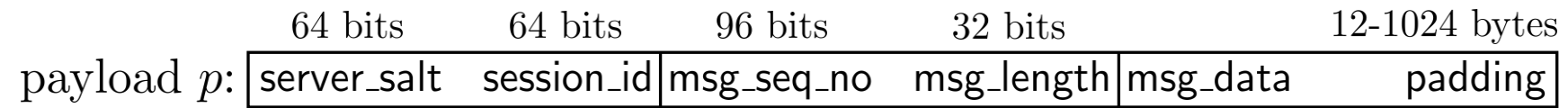
**2018: Kobeissi** reported input validation bugs in Telegram’s Windows Phone client.

**2020: Miculan and Vitacolonna** proved MTProto 2.0 secure in a symbolic model, assuming ideal building blocks.

The focus in the literature has been on the **Secret Chats**.

We focus on the security of the **Cloud Chats**.

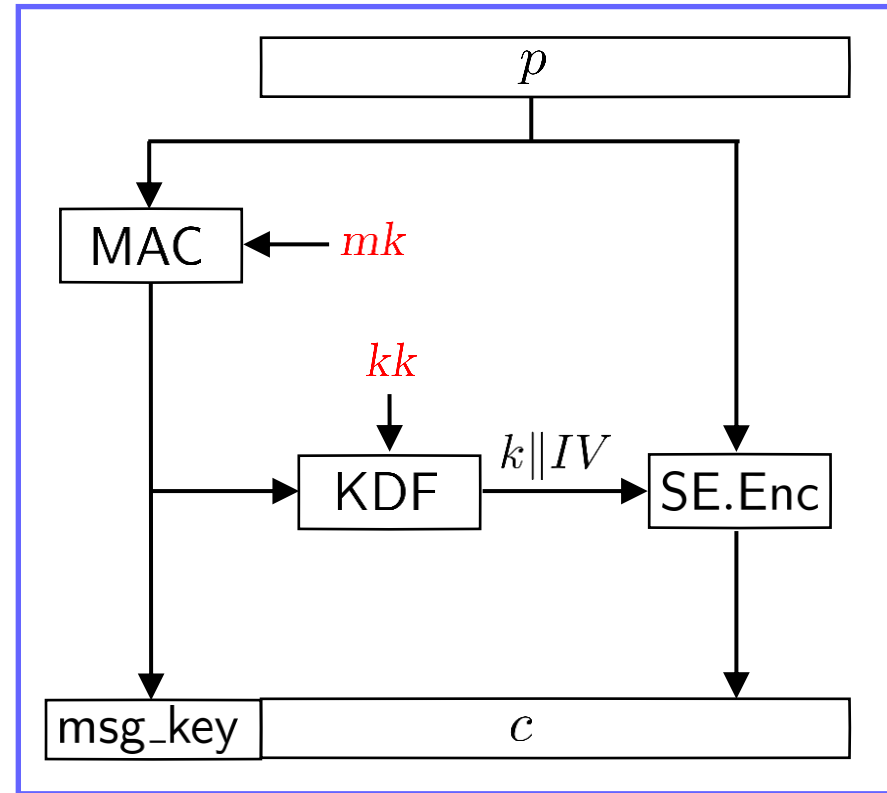
# The Design of MTProto 2.0



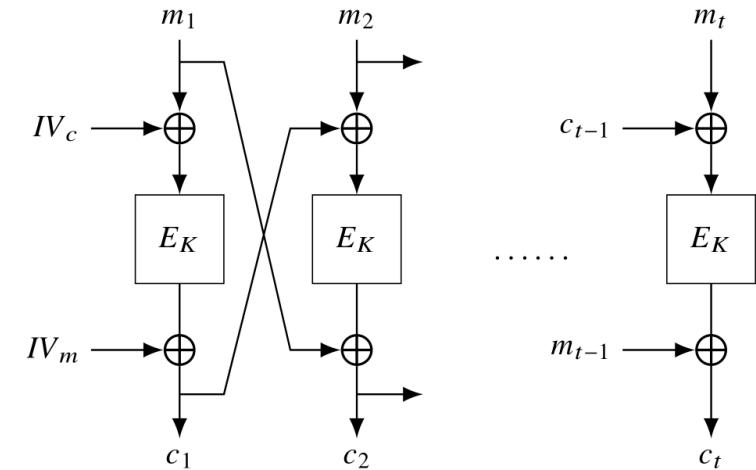
$\text{MAC}(mk, p)$   
 $\text{msg\_key} \leftarrow \text{SHA-256}(mk || p)[64 : 192]$   
Return  $\text{msg\_key}$

$\text{KDF}(kk, \text{msg\_key})$   
 $(kk_0, kk_1) \leftarrow kk$   
 $k_0 \leftarrow \text{SHA-256}(\text{msg\_key} || kk_0)$   
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 $k \leftarrow k_0 || k_1$ ; Return  $k$

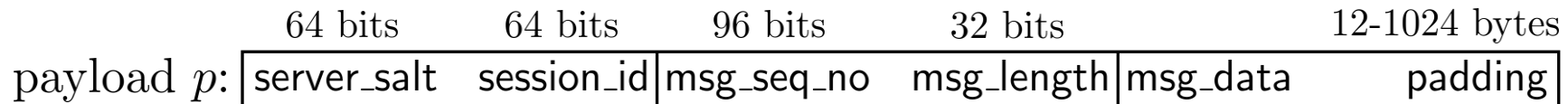
## MTProto.ENCRIPT



## Infinite Garble Extension (IGE) mode:



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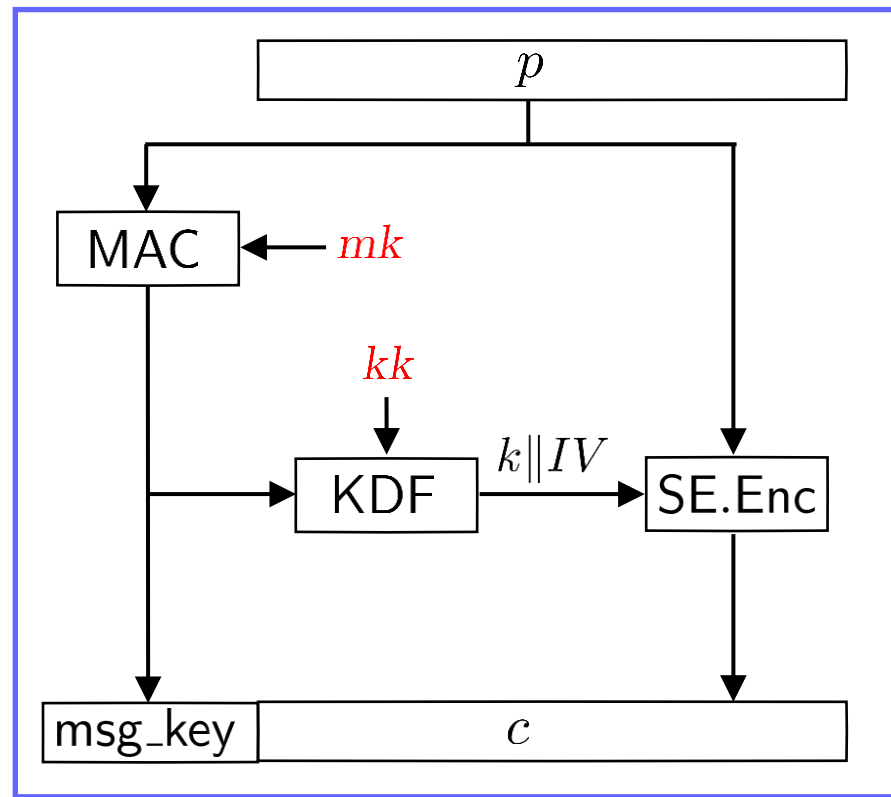
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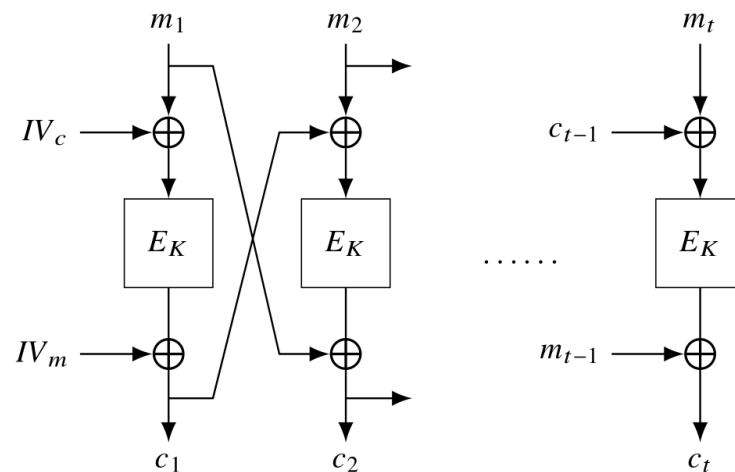
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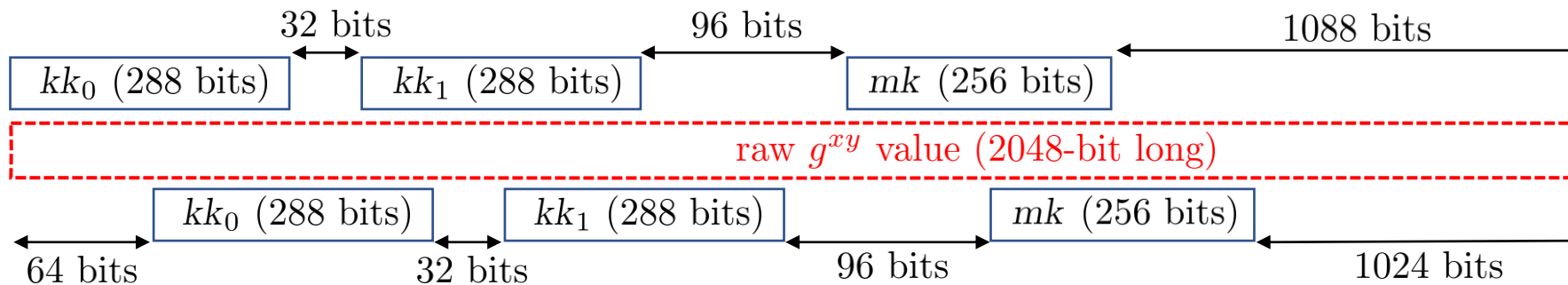


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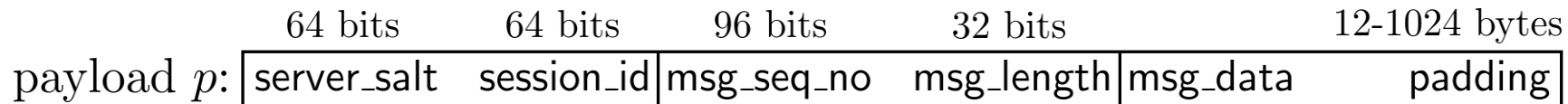


Used to encrypt messages from client to server.

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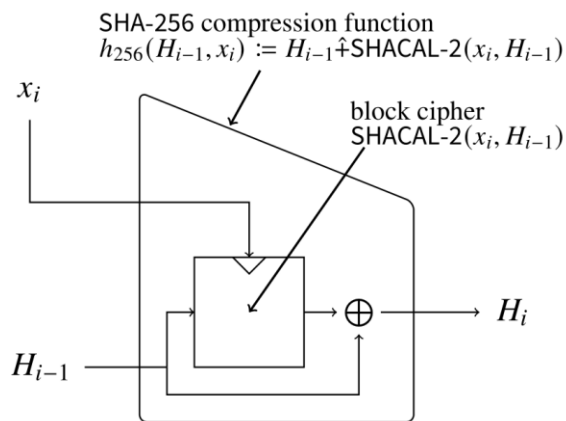
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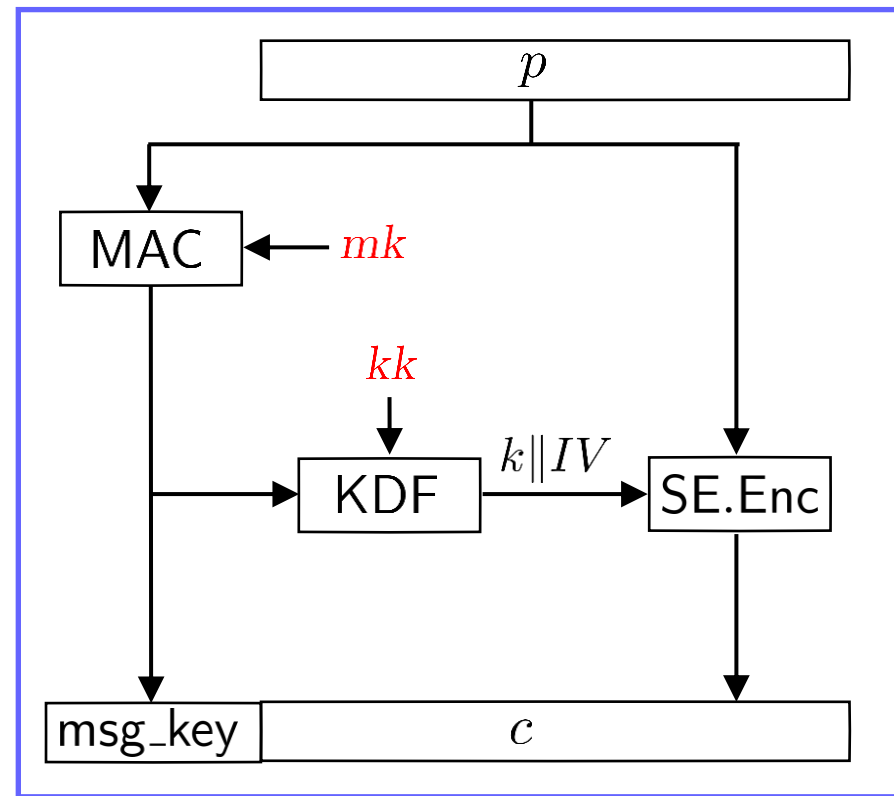
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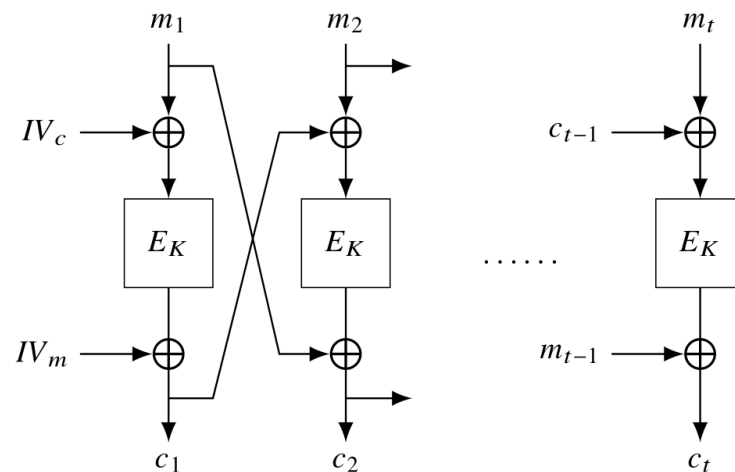
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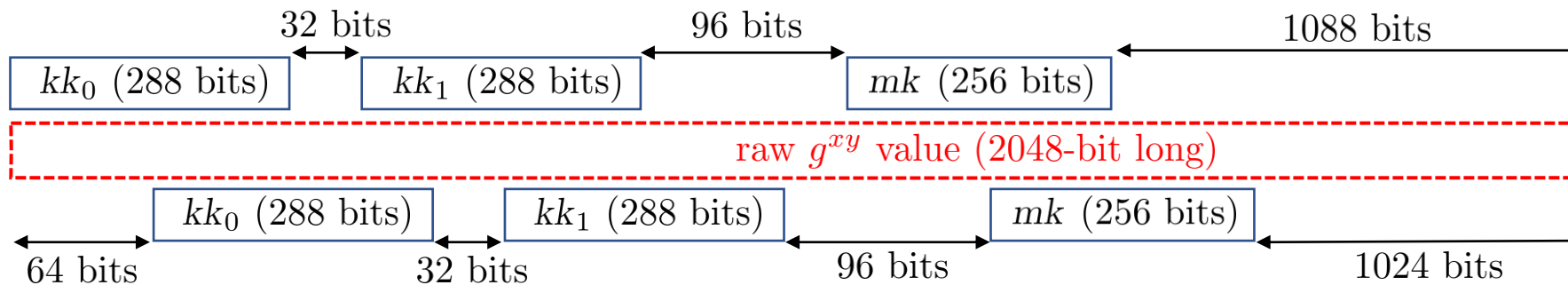


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# Four Attacks Against Telegram

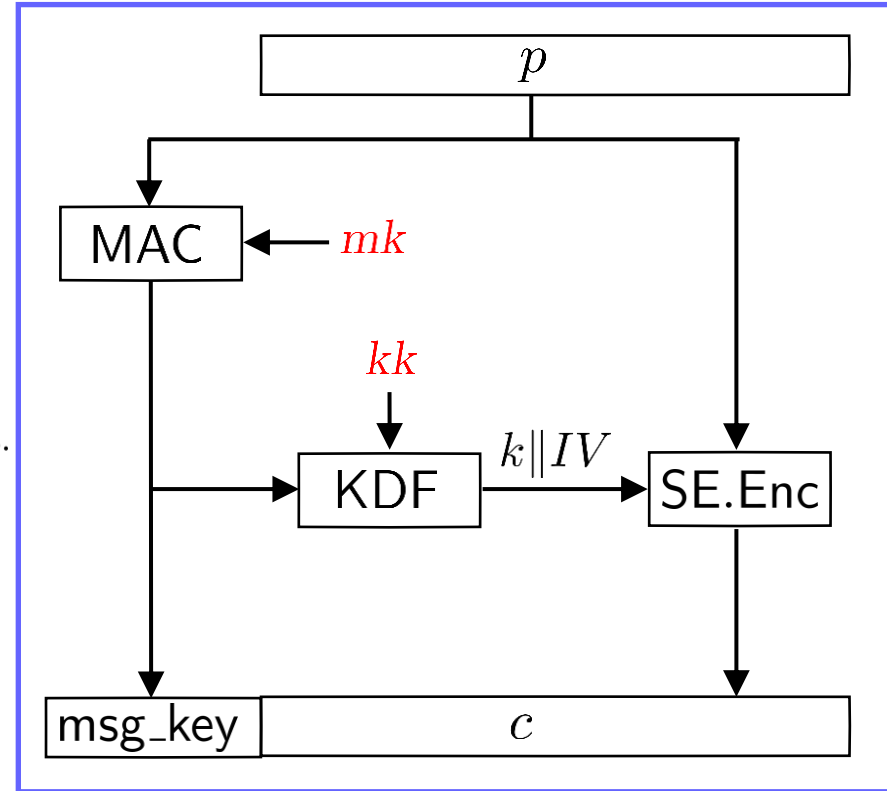
64 bits	64 bits	96 bits	32 bits	12-1024 bytes
server_salt	session_id	msg_seq_no	msg_length	msg_data
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1. Attack against IND-CPA security. // Theoretical.
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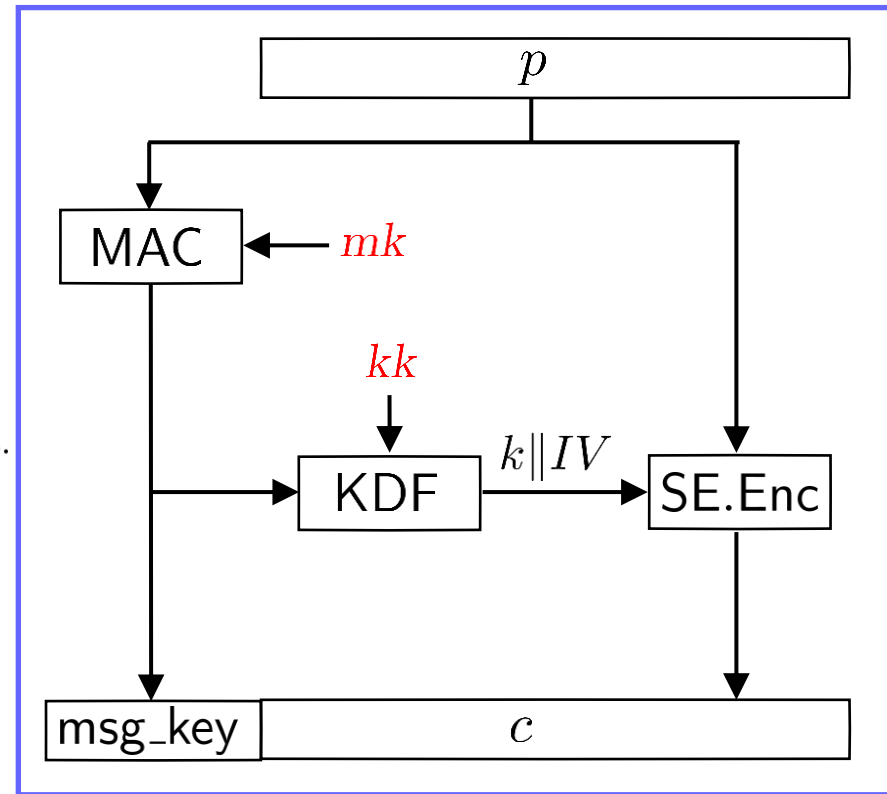
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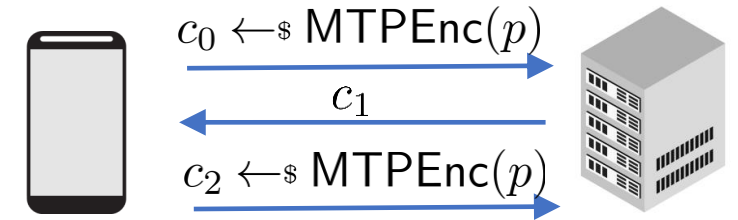
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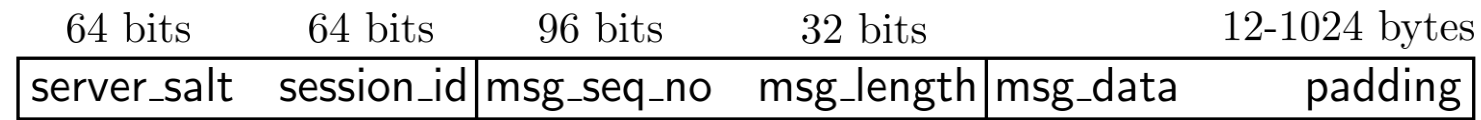
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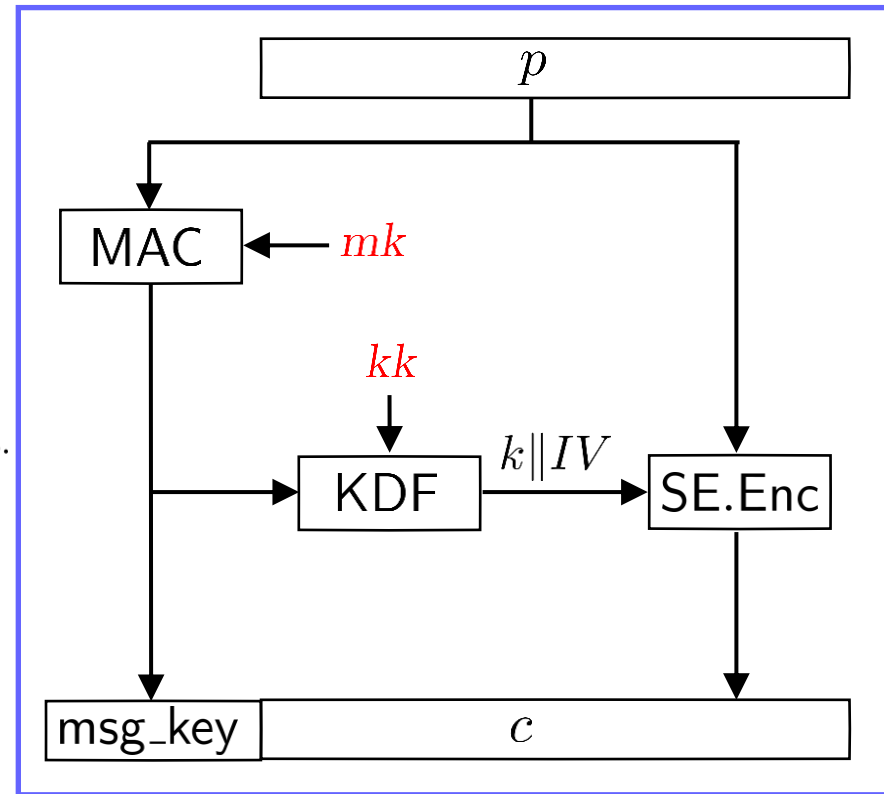


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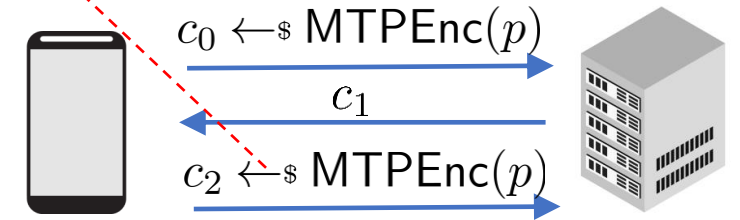
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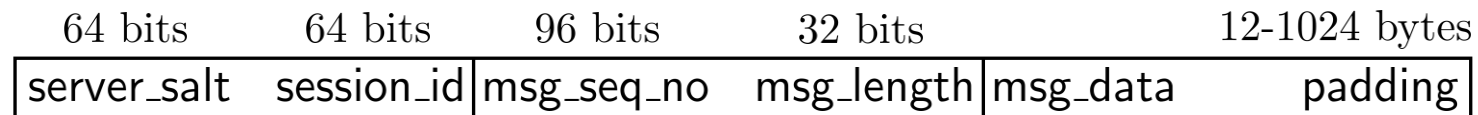


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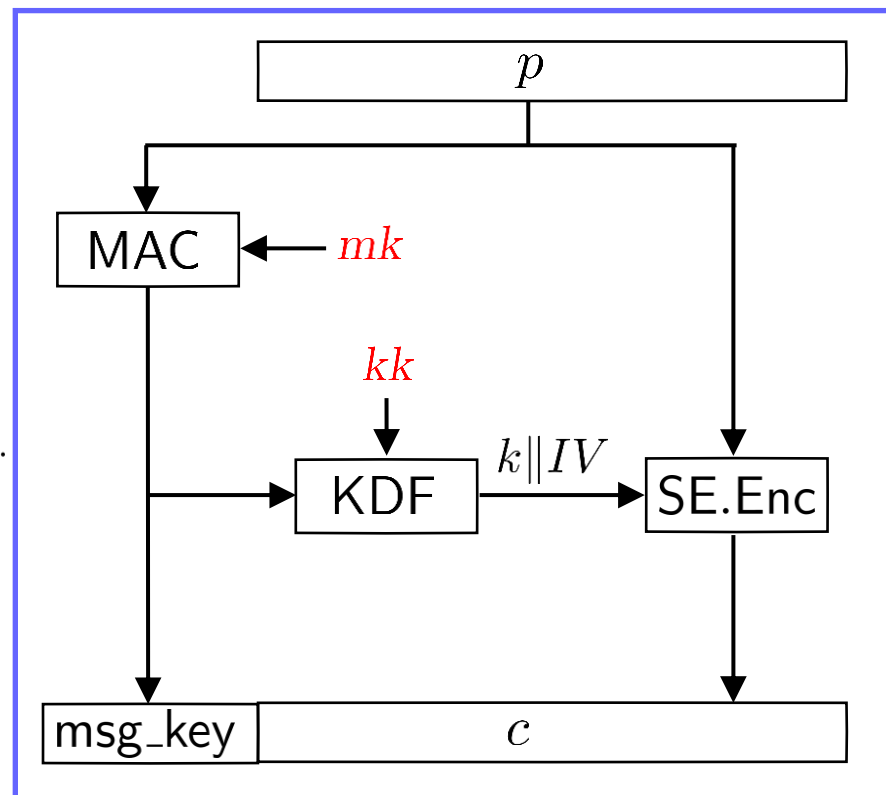


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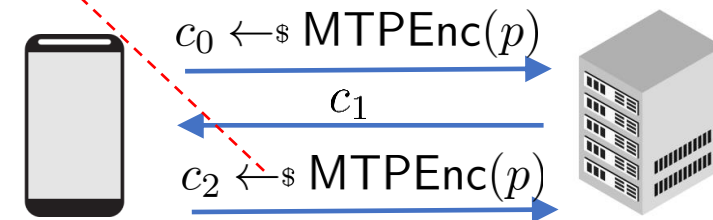
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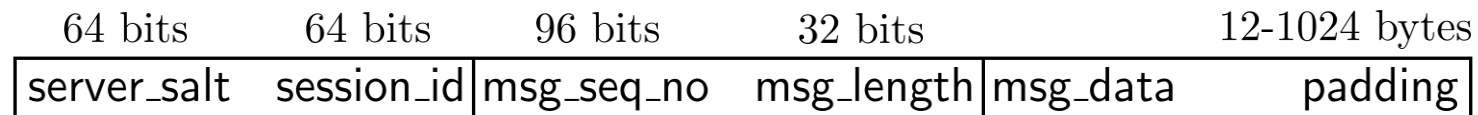


Client expects  $c_1$  be an encryption of ACK. Otherwise, it re-encrypts the payload. Birthday-bound collision in  $msg\_key$  causes the first 2 blocks of  $c_0, c_2$  be the same. IND-CPA thus breaks privacy of  $c_1$ .

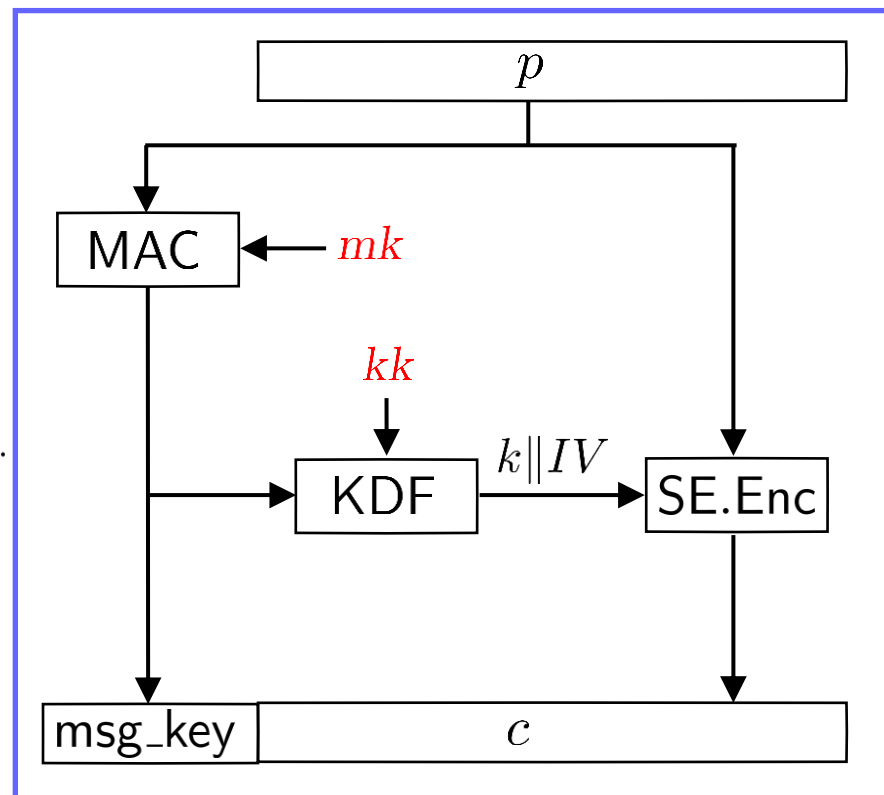
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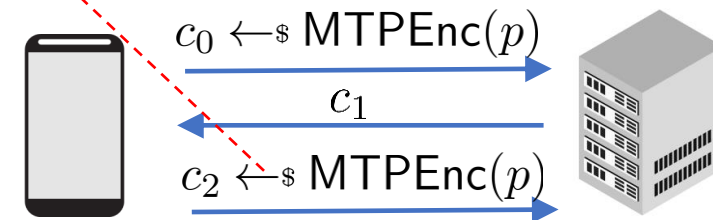


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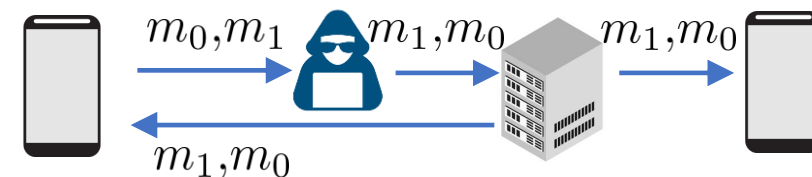


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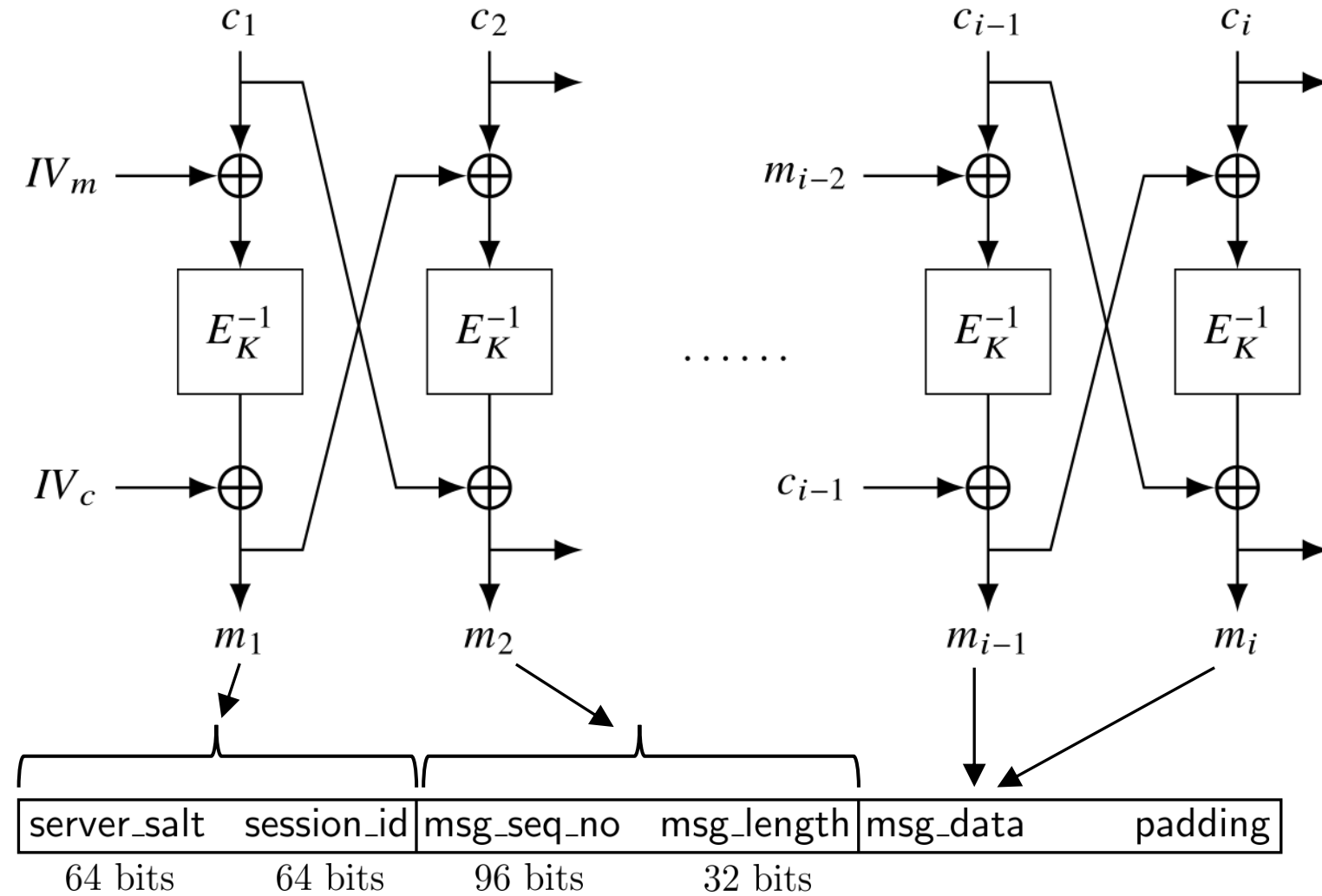
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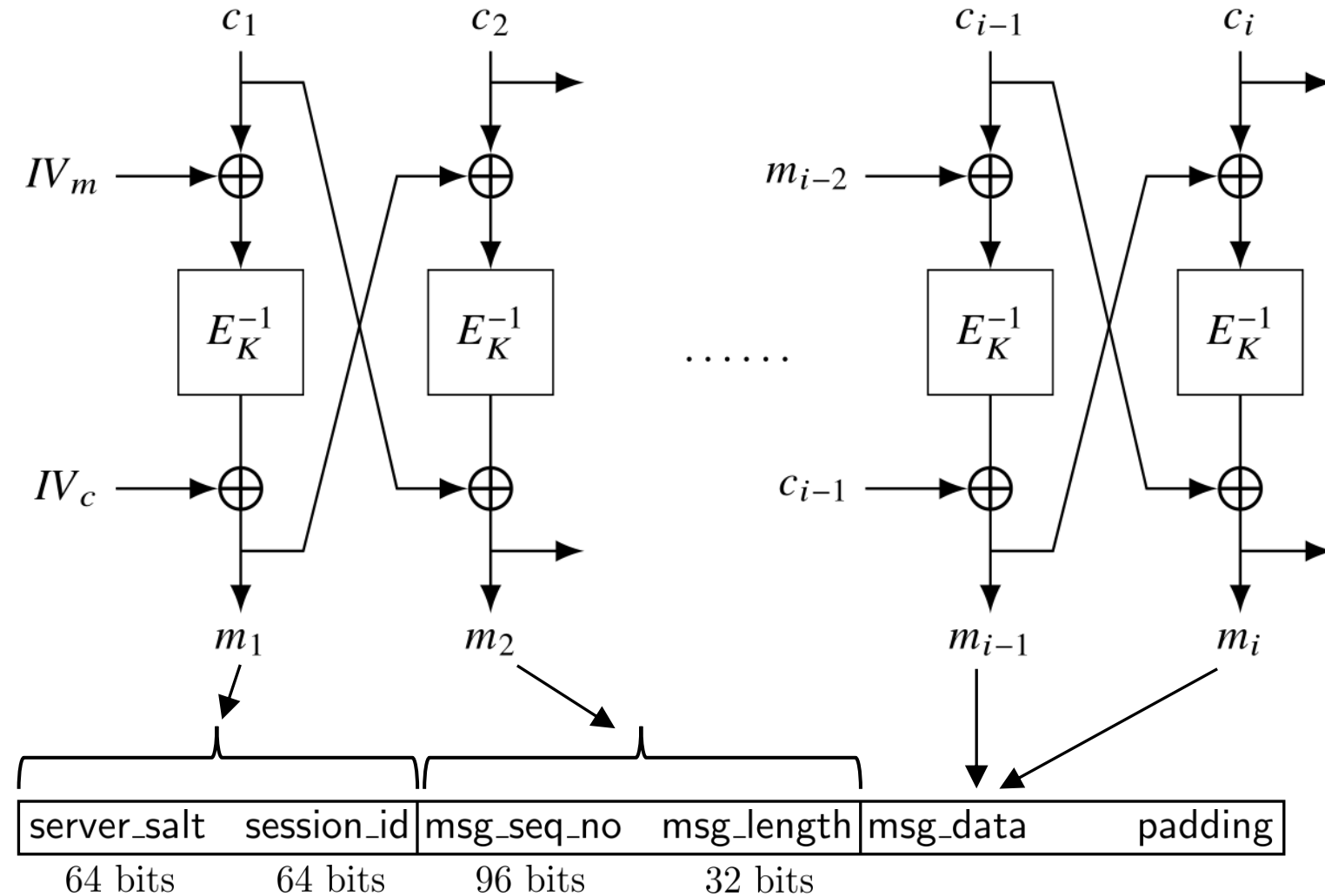
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3 official clients (Android; Desktop; iOS)  
ran a sanity check on a decrypted payload  
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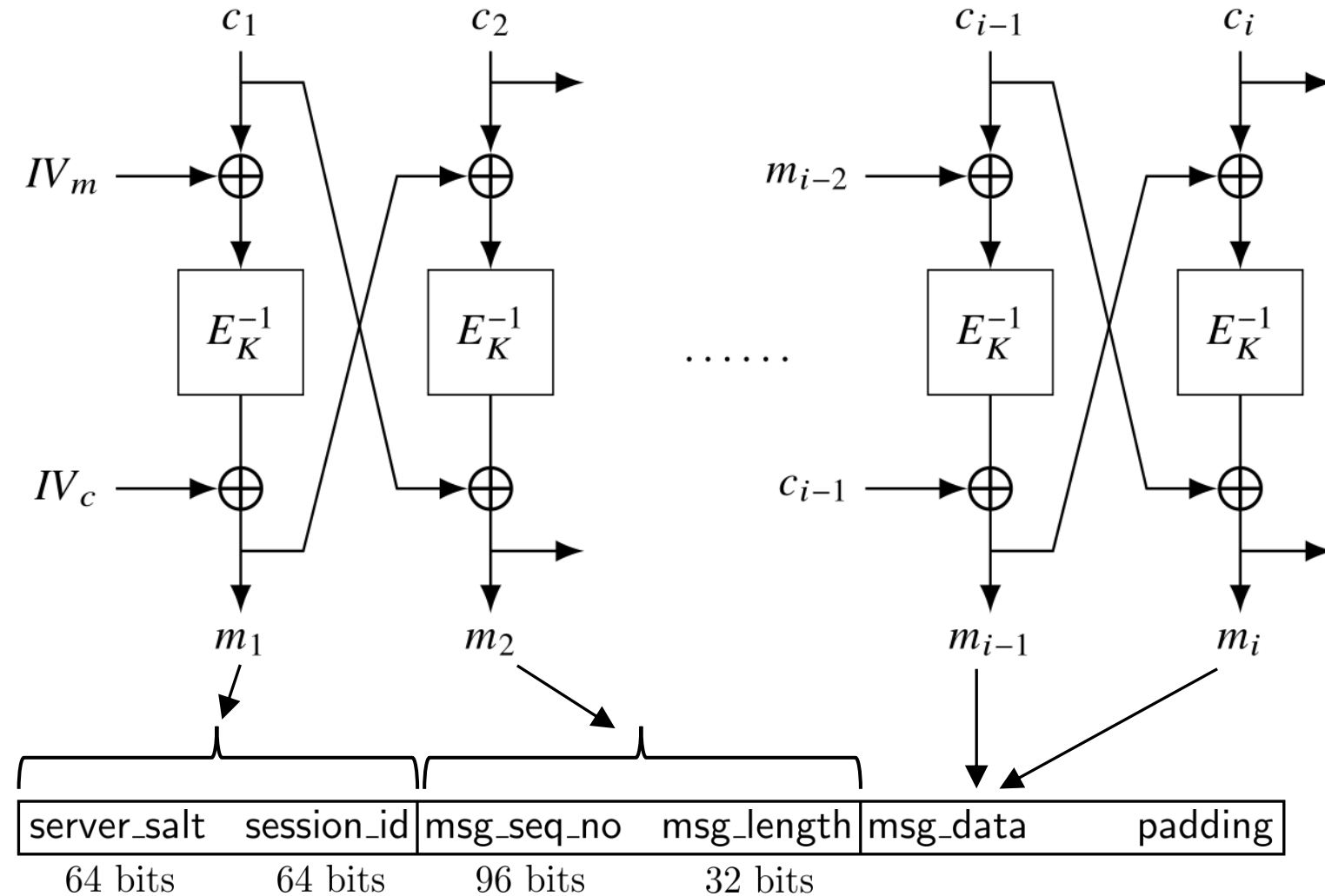
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If (`msg_length > length`) then ... // **Android**  
Outcome of comparison depends on **32 bits** on `msg_length`.  
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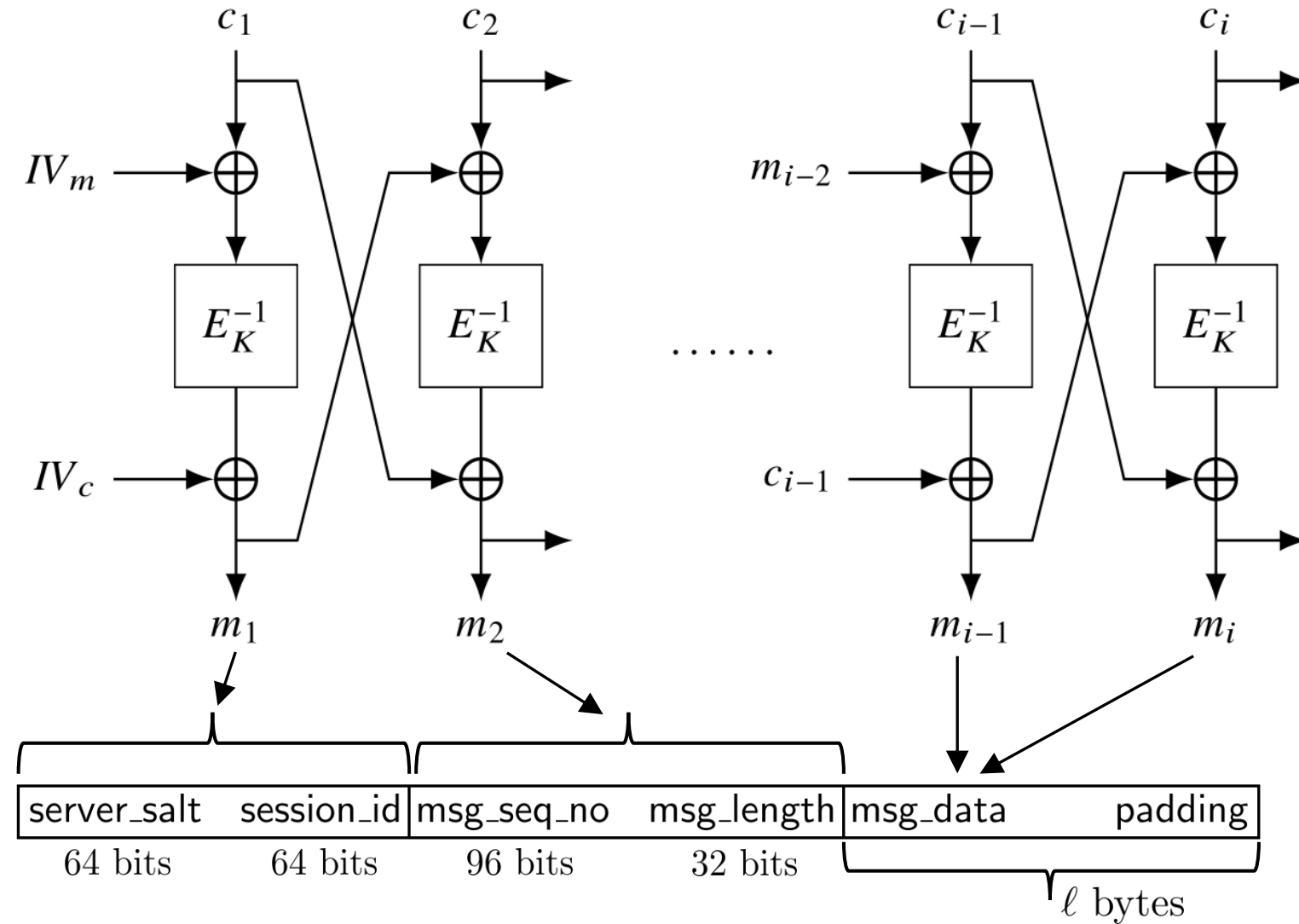
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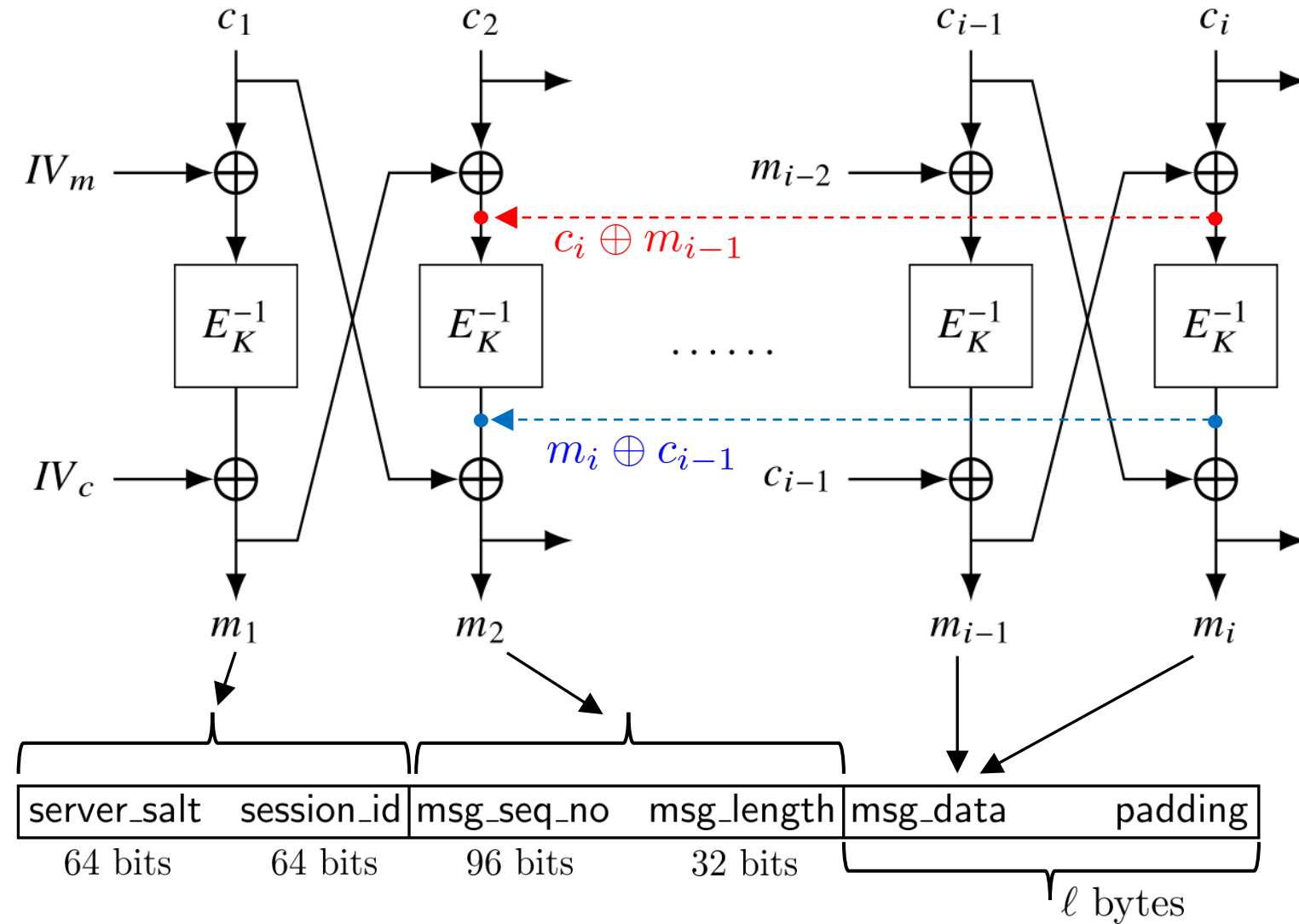
If ( $msg\_length > length$ ) then ... // **Android**  
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If not ( $12 \leq \ell - msg\_length \leq 1024$ ) then ... // **iOS**  
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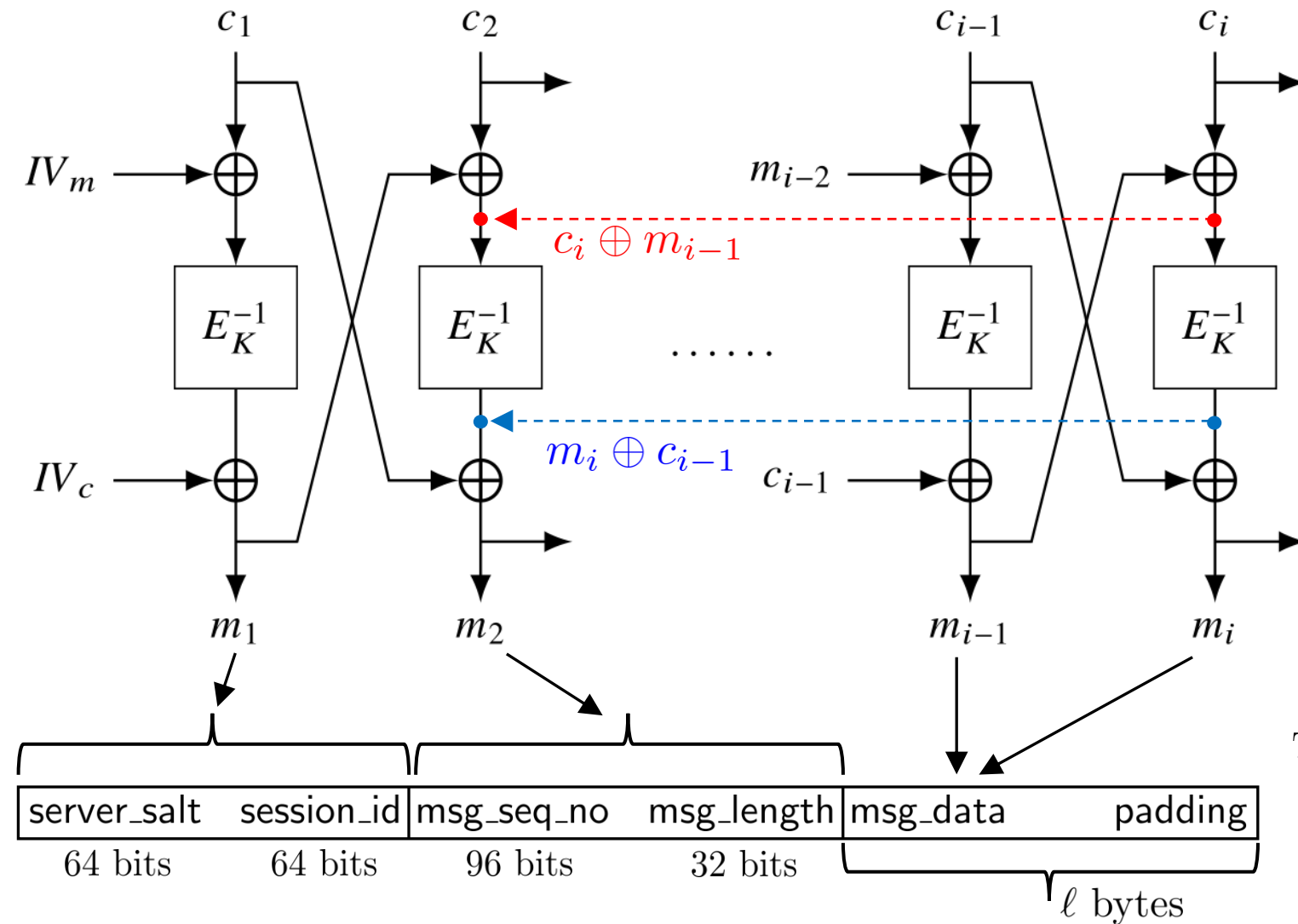
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Assume we know contents of  $m_1$  and  $m_{i-1}$ .  
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 Infer bits of  $m_2$  from timing side-channel.  
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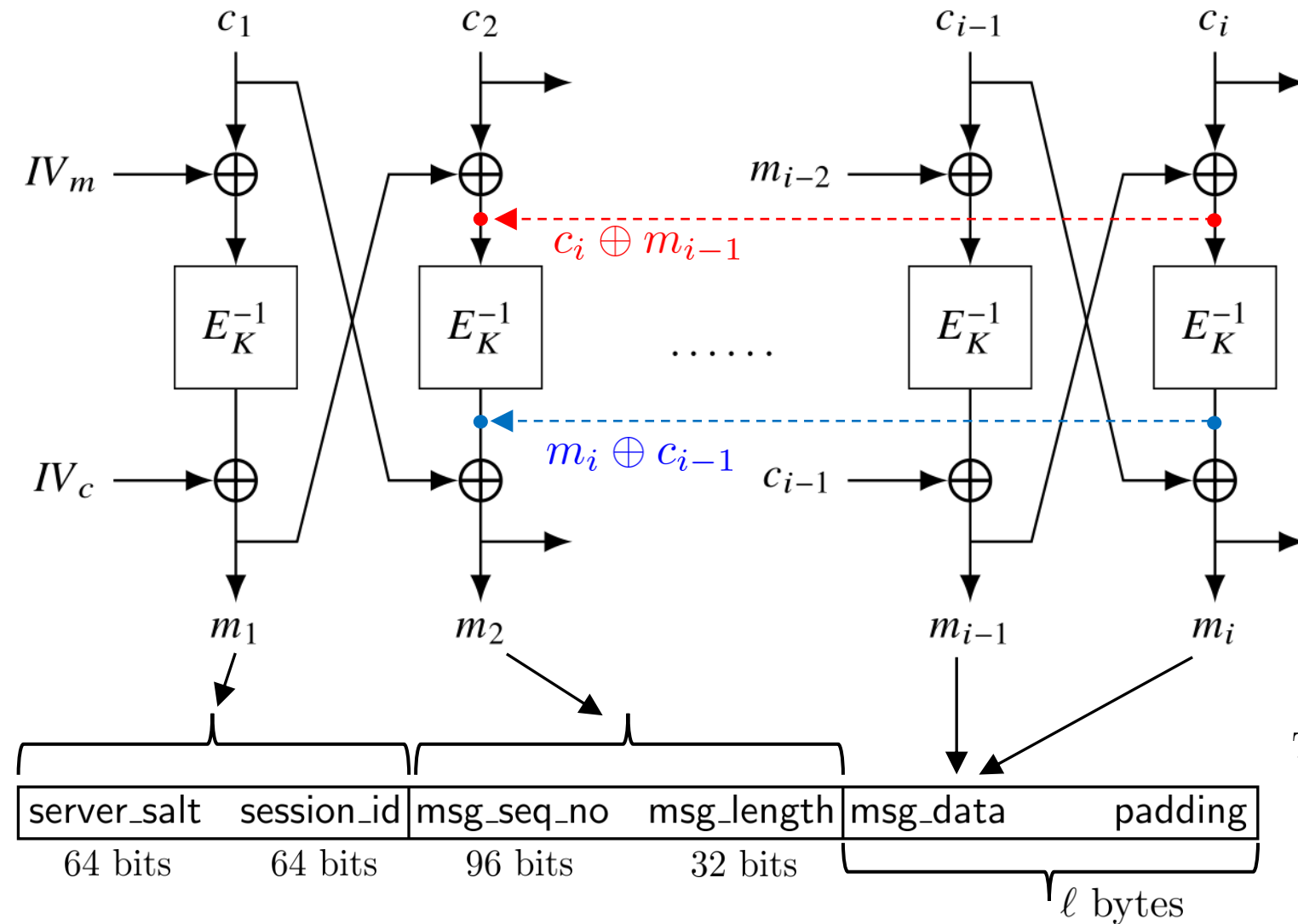
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The attack **fails** because  $server\_salt, session\_id$  are **secret**.

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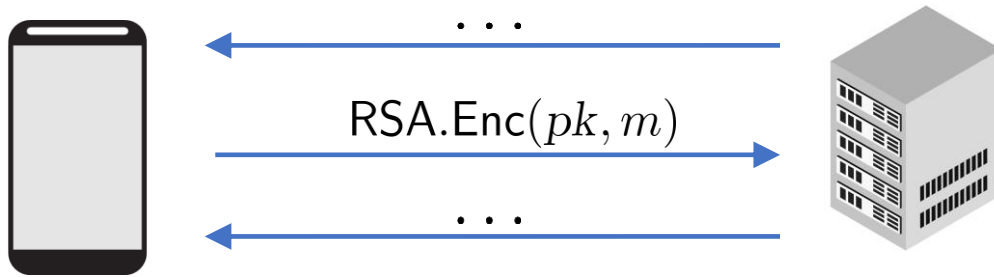
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Our attack highlights the **brittle design**.  
 Stems from using Encrypt-and-MAC.  
 Operates with a decryption key on untrusted data.  
 Would be safer to protect integrity of ciphertext.

# Timing Side-Channel Attack Against Servers

We attack **Telegram**'s key exchange.

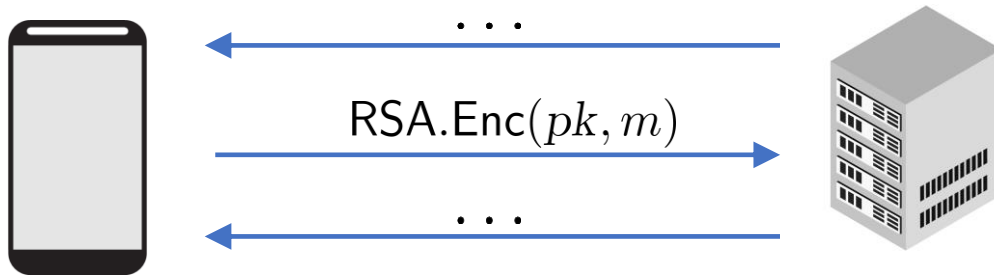


**Telegram** uses textbook **RSA** encryption.

$$m := \text{SHA-1}(\text{data}) \parallel \text{data} \parallel \text{padding}$$

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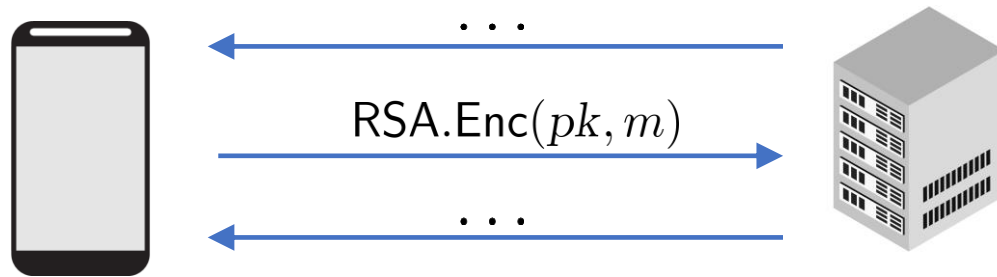
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After RSA decryption but **before** SHA-1 verification, Telegram parses  $m$  to validate its format and might omit the computation of SHA-1.

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**Telegram** uses textbook **RSA** encryption.

$$m := \text{SHA-1}(\text{data}) \parallel \text{data} \parallel \text{padding}$$

After RSA decryption but **before** SHA-1 verification, Telegram parses  $m$  to validate its format and might omit the computation of SHA-1.

We recover data by solving noisy linear equations via **lattice reduction**.

Can use data to recover `server_salt` and `session_id`.

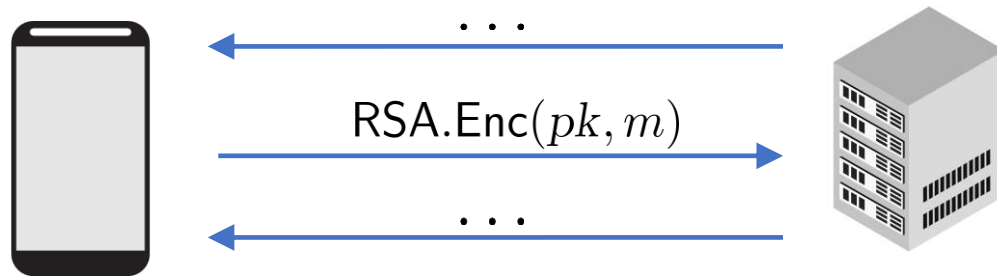
Can use data to run a **MitM** attack against the (encrypted) **DH** exchange.

The attack is **infeasible** in practice because:

- The timing side-channel is very small.
- Recovering `session_id` requires additional  $2^{64}$  computation.
- The key exchange would time out before **MitM** can be completed.

# Timing Side-Channel Attack Against Servers

We attack **Telegram**'s key exchange.



**Telegram** uses textbook **RSA** encryption.

$$m := \text{SHA-1}(\text{data}) \parallel \text{data} \parallel \text{padding}$$

“Publishing the server code doesn’t guarantee privacy, because - unlike with the client-side code - there’s no way to verify that the same code is run on the servers. [..]”

So why not publish the server code anyway, even if it is only a publicity stunt? 3 years ago I learnt that an authoritarian regime [..] was looking for a way to obtain Telegram’s server code. Their plan was to launch their own equally convenient local app and then to shut down all other social media in the country.”

**Pavel Durov** (<https://t.me/durovschat/515221>)

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# Future Work

Large parts of **Telegram**'s design remain unstudied:

Secret chats (including encrypted voice and video calls).

The key exchange.

Multi-user security.

Forward secrecy.

Telegram Passport.

Bot APIs.

The higher-level message processing.

Control messages.

Encrypted CDNs.

Cloud storage.

These are pressing topics for future work.



**Thanks!**

More information at:

<https://mtpsym.github.io/>